

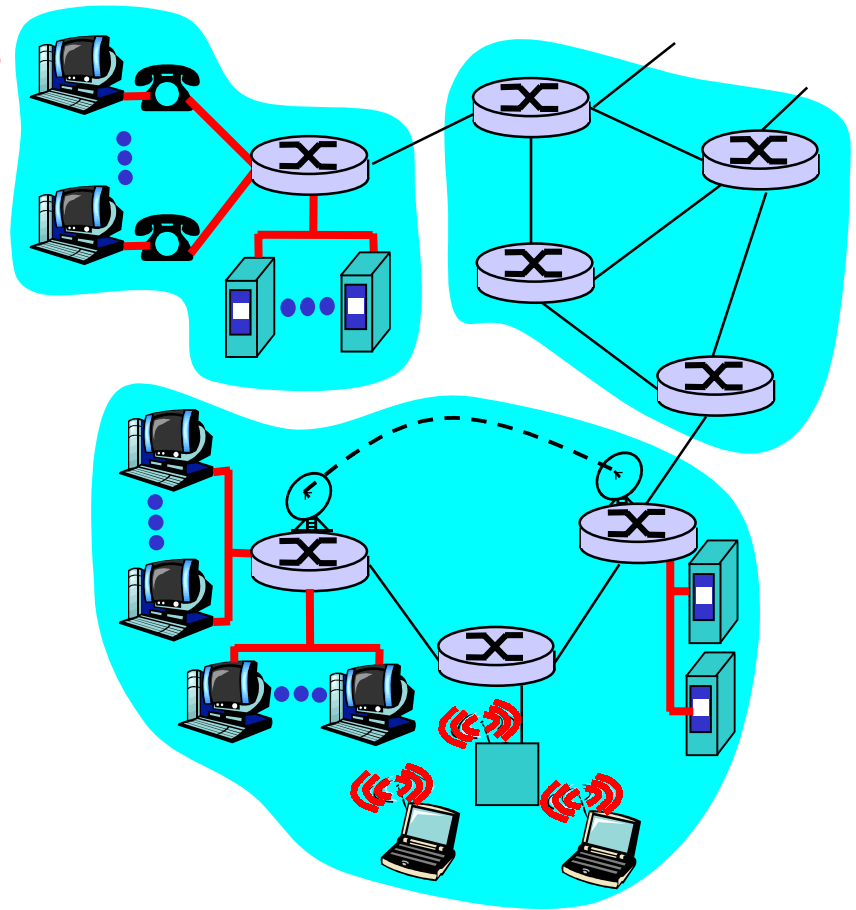
# Access networks and physical media

*Q: How to connect an end system to an edge router?*

- ❑ Residential access nets
- ❑ Institutional access networks (school, company)
- ❑ Mobile access networks

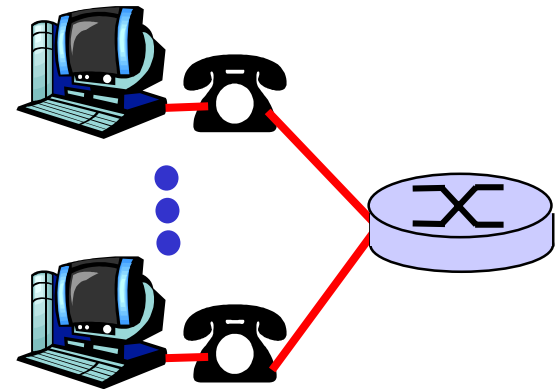
*Keep in mind:*

- ❑ Bandwidth (bits per second) of access network?
- ❑ Shared or dedicated?



# Residential access: point to point access

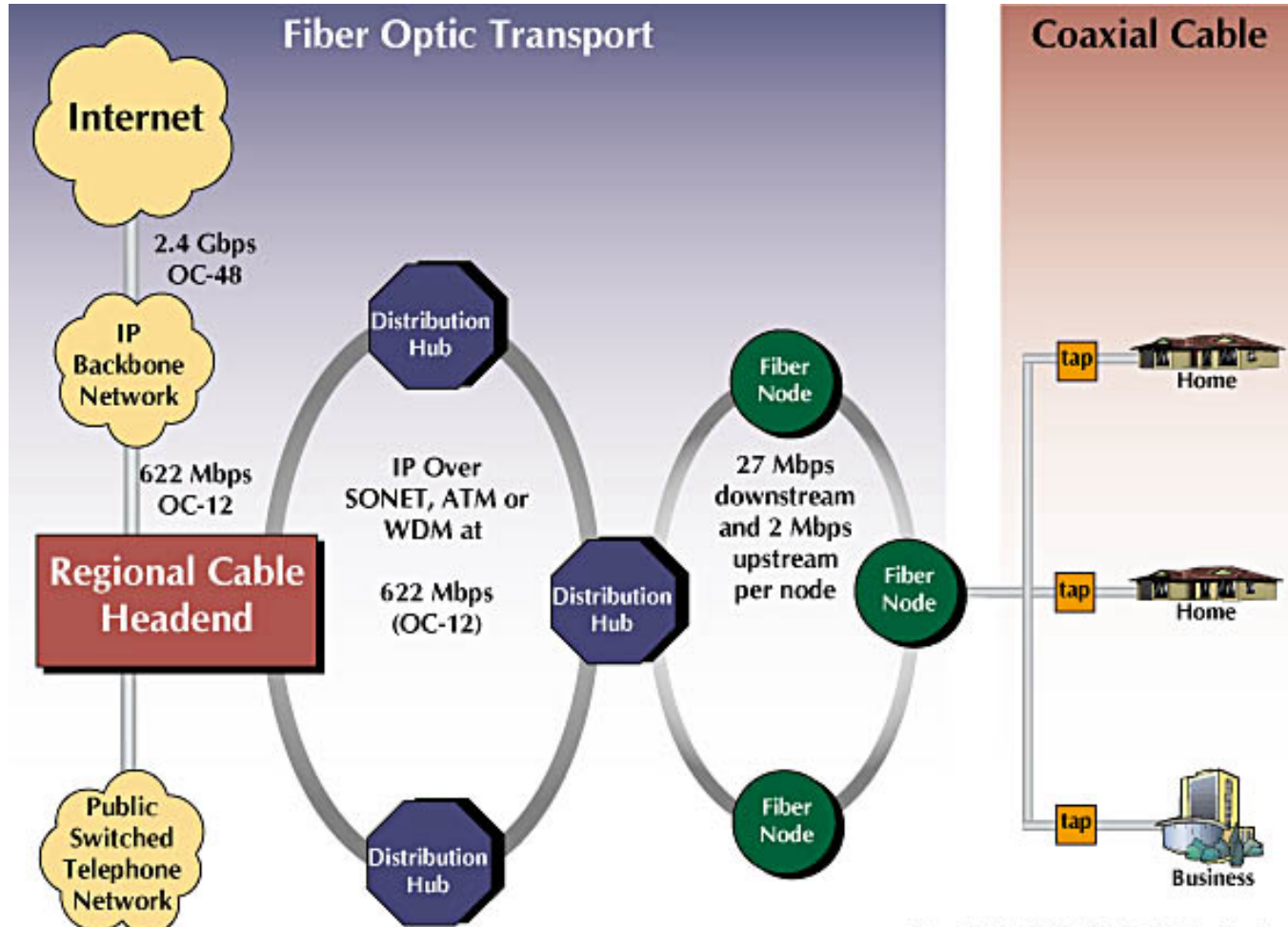
- **Dialup via modem**
  - up to 56Kbps direct access to router (conceptually)
- **ISDN**: integrated services digital network: 128 Kbps all-digital connect to router
- **ADSL**: asymmetric digital subscriber line
  - up to 1 Mbps home-to-router
  - up to 8 Mbps router-to-home
  - ADSL deployment: **happening**



# Residential access: cable modems (1)

- **HFC: hybrid fibre coax**
  - asymmetric: up to 10 Mbps upstream, 1 Mbps downstream
- **Network** of cable and fibre attaches homes to ISP router
  - shared access to router among homes
  - issues: congestion, dimensioning
- **Deployment:** available via cable companies, e.g. MediaOne

# Residential access: cable modems (2)

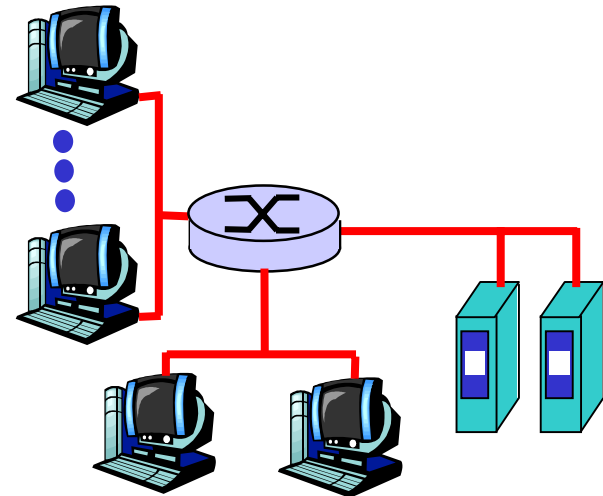


Copyright © 1999 Kinetic Strategies, Inc.

Diagram: <http://www.cabledatcomnews.com/cmhc/diagram.html>

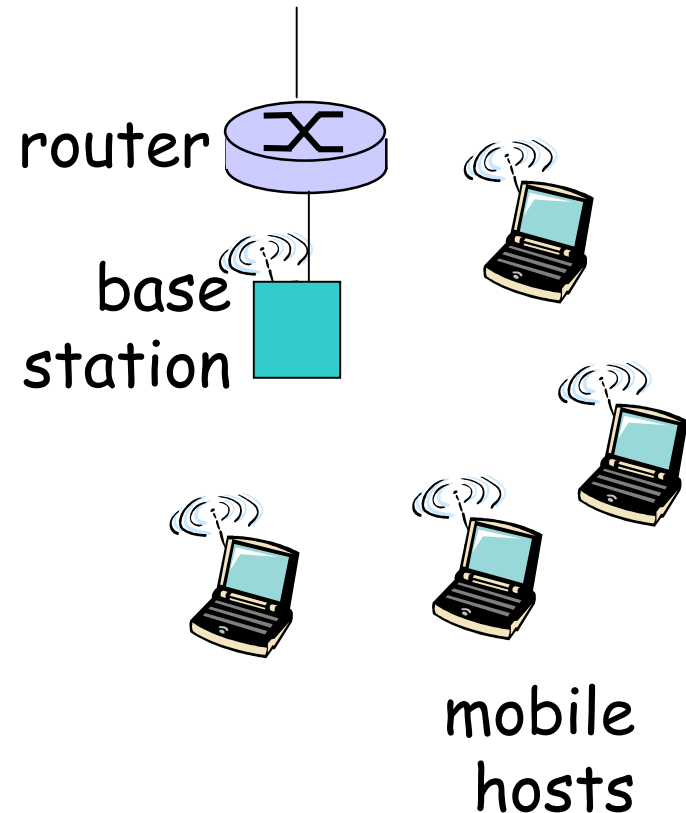
# Institutional access: local area networks

- ❑ Company/univ **local area network** (LAN) connects end system to edge router
- ❑ **Ethernet:**
  - shared or dedicated cable connects end system and router
  - 10 Mbs, 100 Mbps, Gigabit Ethernet
- ❑ **Deployment:** institutions, home LANs happening now
- ❑ LANs: chapter 5



# Wireless access networks

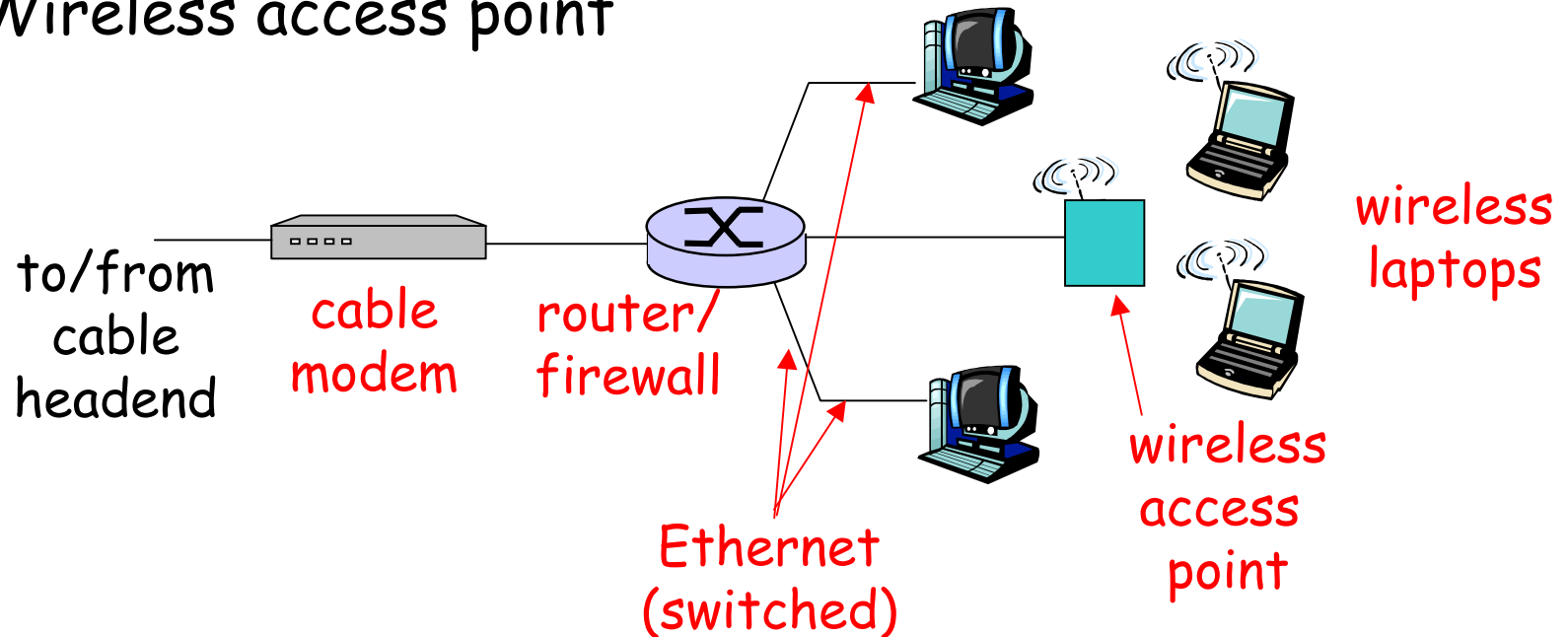
- ❑ Shared *wireless* access network connects end system to router
- ❑ **Wireless LANs:**
  - radio spectrum replaces wire
  - e.g., Lucent Wavelan 11 Mbps
- ❑ **Wider-area wireless access**
  - CDPD: wireless access to ISP router via cellular network



# Home networks

## Typical home network components:

- ❑ ADSL or cable modem
- ❑ Router/firewall
- ❑ Ethernet
- ❑ Wireless access point



# Physical media

- **Physical link:**  
transmitted data bit propagates across link
- **Guided media:**
  - signals propagate in solid media: copper, fibre
- **Unguided media:**
  - signals propagate freely, e.g., radio

## Twisted Pair (TP)

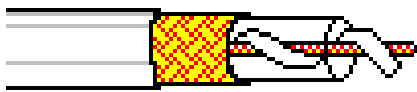
- Two insulated copper wires
  - Category 3: traditional phone wires, 10 Mbps Ethernet
  - Category 5 TP: 100 Mbps Ethernet



# Physical media: coax, fibre

## Coaxial cable:

- ❑ Wire (signal carrier) within a wire (shield)
  - baseband: single channel on cable
  - broadband: multiple channel on cable
- ❑ Bidirectional
- ❑ Common use in 10 Mbs Ethernet



## Fibre optic cable:

- ❑ Glass fibre carrying light pulses
- ❑ High-speed operation:
  - 100 Mbps Ethernet
  - high-speed point-to-point transmission (e.g., 5 Gps)
- ❑ Low error rate



# Physical media: radio

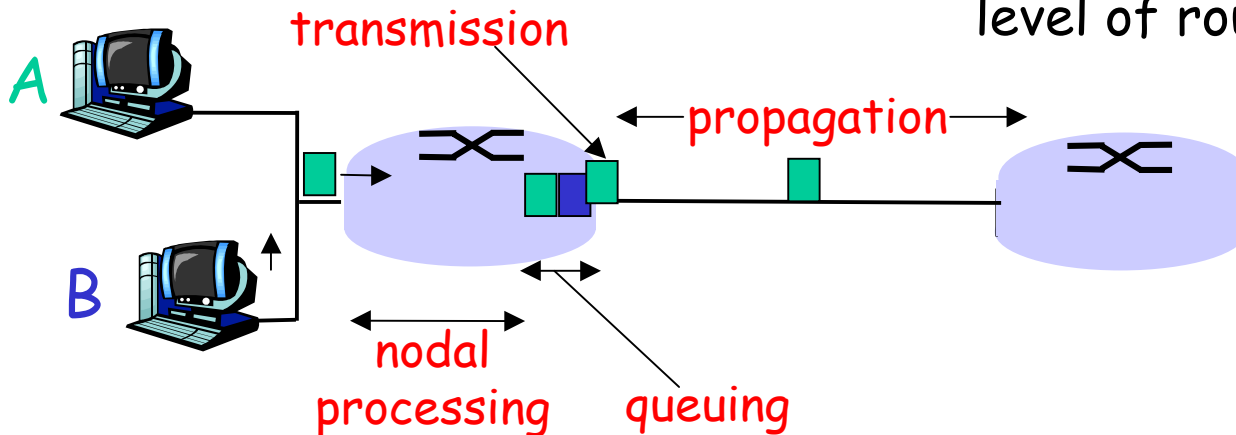
- ❑ Signal carried in electromagnetic spectrum
  - ❑ No physical 'wire'
  - ❑ Bidirectional
  - ❑ Propagation
- Environment effects:
- reflection
  - obstruction by objects
  - interference

## Radio link types:

- ❑ **Microwave**
  - e.g. up to 45 Mbps channels
- ❑ **LAN** (e.g., WaveLAN)
  - 2 Mbps, 11 Mbps
- ❑ **Wide-area** (e.g. cellular)
  - e.g. CDPD, 10's Kbps
- ❑ **Satellite**
  - up to 50 Mbps channel (or multiple smaller channels)
  - 270 Msec end-end delay
  - geosynchronous versus LEOS

# Delay in packet-switched networks (1)

- Packets experience **delay** on end-to-end path
- **Four** sources of delay at each hop
  - Nodal processing:
    - check bit errors
    - determine output link
  - Queuing
    - time waiting at output link for transmission
    - depends on congestion level of router



# Delay in packet-switched networks (2)

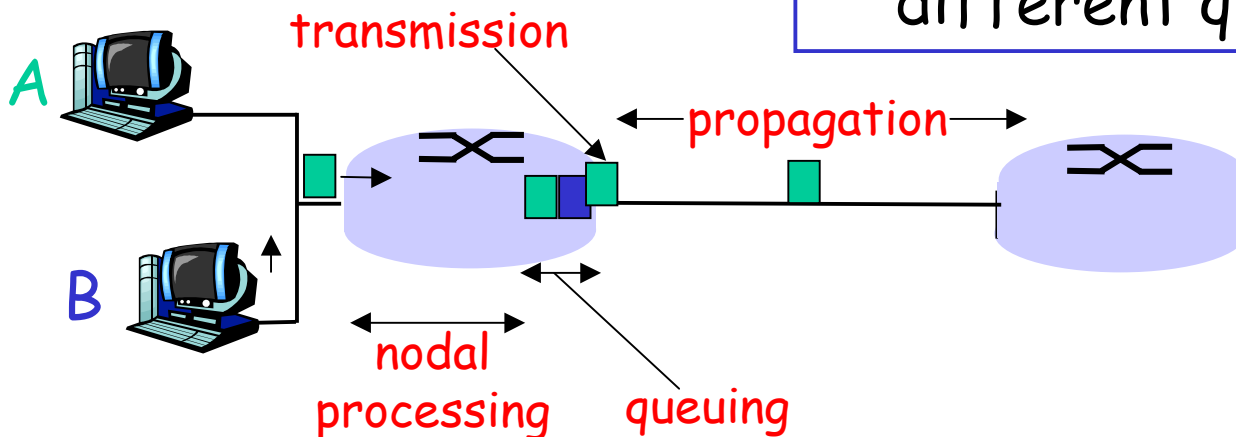
## Transmission delay:

- $R$  = link bandwidth (bps)
- $L$  = packet length (bits)
- time to send bits into link =  $L/R$

## Propagation delay:

- $d$  = length of physical link
- $s$  = propagation speed in medium ( $\sim 2 \times 10^8$  m/sec)
- propagation delay =  $d/s$

**Note:**  $s$  and  $R$  are very different quantities!

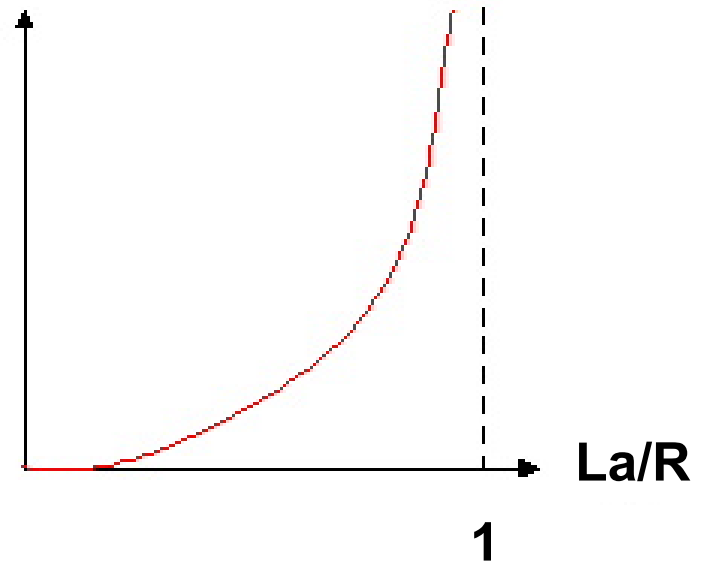


# Queuing delay (revisited)

- $R$ =link bandwidth (bps)
- $L$ =packet length (bits)
- $a$ =average packet arrival rate

traffic intensity =  $La/R$

Average queuing  
delay



- $La/R \sim 0$ : average queuing delay small
- $La/R \rightarrow 1$ : delays become large
- $La/R > 1$ : more 'work' arriving than can be serviced, average delay infinite!

# 'Real' Internet delays and routes

**traceroute:** routers, rt delays on source-dest path  
also: pingplotter, various windows programs

```
1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 * * *
18 * * *
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms
```

# Protocol 'layers'

## Networks are complex!

- Many 'pieces':
  - hosts
  - routers
  - links of various media
  - applications
  - protocols
  - hardware, software

## Question:

Is there any hope of *organizing* the structure of a network?

Or at least our discussion of networks?

# Internet history

## *1961-1972: Early packet-switching principles*

- 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- 1964: Baran - packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: First ARPAnet node operational
- 1972:
  - ARPAnet demonstrated publicly
  - NCP (Network Control Protocol) first host-host protocol
  - First e-mail program
  - ARPAnet has 15 nodes

# Internet history

## *1972-1980: Internetworking, new and proprietary nets*

- ❑ 1970: ALOHAnet satellite network in Hawaii
- ❑ 1973: Metcalfe's PhD thesis proposes Ethernet
- ❑ 1974: Cerf and Kahn - architecture for interconnecting networks
- ❑ late70s: proprietary architectures: DECnet, SNA, XNA
- ❑ late 70s: switching fixed length packets (ATM precursor)
- ❑ 1979: ARPAnet has 200 nodes

### *Cerf and Kahn's internetworking principles:*

- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

*define today's Internet architecture*

# Internet history

## *1980-1990: new protocols, a proliferation of networks*

- ❑ 1983: deployment of TCP/IP
- ❑ 1982: smtp e-mail protocol defined
- ❑ 1983: DNS defined for name-to-IP-address translation
- ❑ 1985: ftp protocol defined
- ❑ 1988: TCP congestion control
- ❑ New national networks: Cset, BITnet, NSFnet, Minitel
- ❑ 100,000 hosts connected to confederation of networks

# Internet history

## *1990s: commercialization, the WWW*

- ❑ Early 1990s: ARPAnet decommissioned
- ❑ 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- ❑ early 1990s: WWW
  - hypertext [Bush 1945, Nelson 1960s]
  - HTML, http: Berners-Lee
  - 1994: Mosaic, later Netscape
  - late 1990s: commercialization of the WWW

## Late 1990s:

- ❑ est. 50 million computers on Internet
- ❑ est. 100 million+ users
- ❑ backbone links running at 1 Gbps

# Summary

## Covered:

Internet overview

- ❑ What is a protocol?
- ❑ Network edge, core, access network
  - packet-switching versus circuit-switching
- ❑ Performance: loss, delay
- ❑ Layering and service models
- ❑ Backbones, NAPs, ISPs
- ❑ History

## You now have:

- ❑ Context, overview, 'feel' of networking
- ❑ More depth, detail *later* in course